

20.05.14

Ann-Algo

Q = prequeue of pointers to nodes in kd-tree
sorted by distance to q (query pt)

P^o = candidate of NN,
init $v = \text{root}$
 $P^o = P^t$ at ∞
 $Q = \text{empty}$

while $d(v, q) < \frac{1}{1+\epsilon} d(P^o, q)$

while v is inner node

let v_1, v_2 be children of v ,
with v_1 closer to q

insert v_2 into Q

$v := v_1$

if $d(p_v, q) < d(P^o, q)$.

$P^o = p_v$

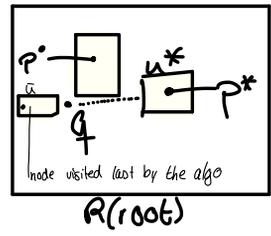
$v = \text{extractMin}(Q)$

... return P^o

Correctness:

Let u^* be leaf containing $p^* = u$

- a) algo visits $u^* \rightarrow p^o = p^*$
- b) u^* is not visited $\rightarrow p^o \neq p^*$



$$d(p^*, q) \geq d(q, u^*) \geq d(q, \bar{u})$$

\downarrow
d is metric
 \downarrow
closer nodes are visited first

$d(q, \bar{u}) \geq \frac{1}{1+\epsilon} d(q, p^o)$ // by condition in the outer while

$$\Rightarrow d(q, p^*) \geq \frac{1}{1+\epsilon} \cdot d(q, p^o) \Rightarrow q^o \text{ is } (1+\epsilon)\text{-ANN}$$

Time

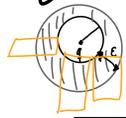
One outer while iteration: 1x extract Min

innerloop. $O(\log n)$ times insert op. into 1

Normal heaps $O(\log n)$ for extract & insert $\Rightarrow ANN \in O(l \cdot \log^2 n)$

Fibonacci heaps: $O(\log n)$ extract, insert oper. has $O(1)$ amortized time. $\Rightarrow ANN \in O(l \cdot \log n)$

$$l = \# \text{leaves} = O(\log^{d-1} n)$$



[argument considers # boxes that can penetrate annulus of thickness ϵ , where boxes emerge from kd-tree with longest-side splitting!]

"Best" ANN algorithm

Randomized kd-tree (RkD):

For a split: - determine D many axes (dim's) with the longest variance

- choose one of those randomly

- split across median

$D \approx 5$ is fine in most cases

RkD forest.

build several RkD's over pt set P

ANN search:

maintain one p -queue (for all RkD's forest)

beginning: descent into each RkD tree down to the leaf containing q (= first iteration of outer while in orig)

- put the "other" nodes (vz's) along the way into Q



proceed as before with orig ANN algo

Application: texture synthesis

Input: Image I , texture, i.e., color of pixel depends on neighborhood only

Output: Image T , bigger, looking "similar" to I

Notation: $p_i := \text{pixel} \in I$

$p_o := \text{output pixel} \in T$

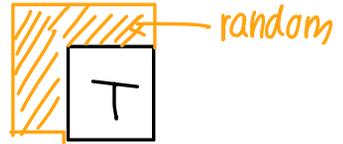
$\mathcal{N}(p) = \text{neighborhood around } p$
(shape depends on specifics of algo)



Algorithm:

Init $T := \text{empty}$, plus random border around randomly colors

for all p_o 's $\in T$ in scan line order:



NN search with d -dim pts, where $d = \# \text{pixels in } \mathcal{N}(p)$

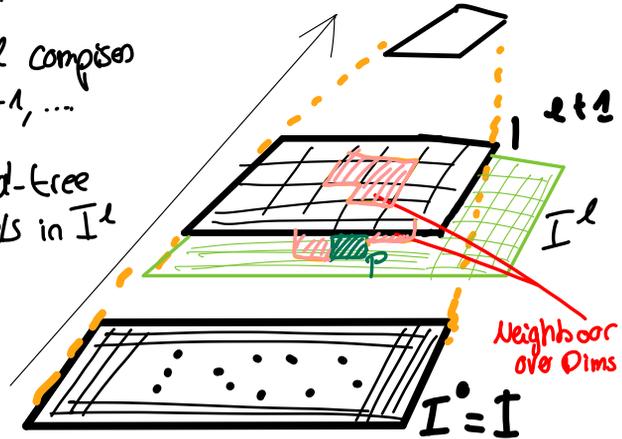
$\left\{ \begin{array}{l} \text{find } p_i \in I \text{ with} \\ d(\mathcal{N}(p_o), \mathcal{N}(p_i)) = \min \\ p_o = p_i \end{array} \right.$

Build image pyramid over I

$I_0 = I$, I^{l+1} created from I^l
by averaging 2×2 pixels (or other smoothing op)

Neighborhood $N(p)$ on level l comprises
pixels on levels $l, l+1, \dots$

For each level l : build kd-tree
over all $N(p)$'s for all pixels in I^l



Build image pyramid for T , but top-down!

On every level l : images T^{l+1}, \dots
are already created,
create random border around T^l ,
proceed as before.